

Dyckhoff

# Secondary Dold-Kan

①

① "Pre-primary" integers

① "Primary" ab gps (Dold-Kan as usual)

② "Secondary" dg cats or stable cocat.

} category

It is Kapranov  
 Serre  
 Suhlmann  
 Jasso (anahitaii c.c.3)  
 Poguntke

①  $\mathcal{X}$  eiter char

①  $H_0(\ ; A)$  homdy.

② 2ndary homdy  
 w/ cat in  $\mathcal{C}_i$   
 compdy file in  
 top. setty.  
 "Borel-Moore"  $\leftrightarrow$  wrapped filey.

① Preprimary:  $a = (a_0, a_1, a_2, \dots)$  cequence of integers "discrete fun"

$\rightsquigarrow$  cequence of "forward differences"

$\Delta a = (a_1 - a_0, a_2 - a_1, a_3 - a_2, \dots)$  "discrete derivate"

iterate  $\rightsquigarrow$  Newton's formula

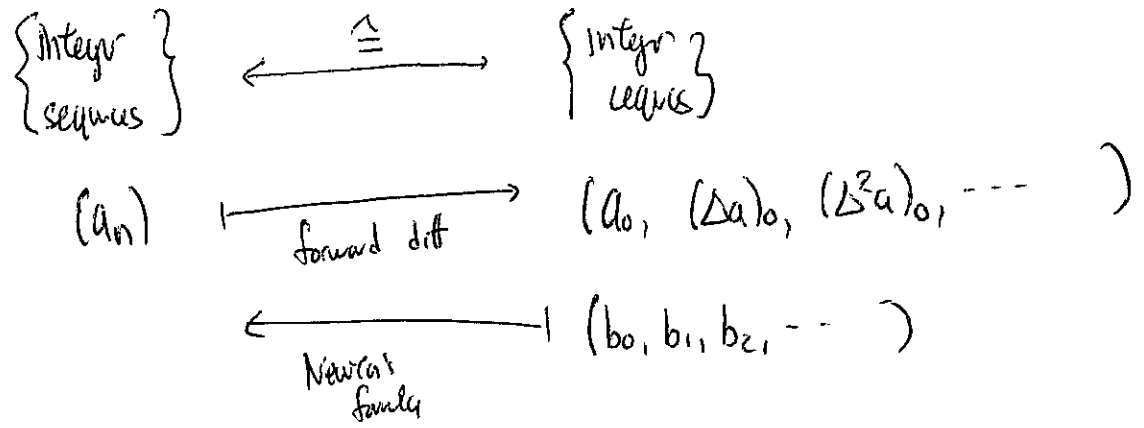
$$a_n = \sum_{k=0}^n \binom{n}{k} (\Delta^k a)_0$$

"discrete Taylor's formula"

$$f(x) = \sum \frac{x^k}{k!} f^{(k)}(0)$$

Organize to a correspondence:

(2)



Ex Computing  $\Delta^m a$  can tell you whether  $(a_i)$  is polynomial: i.e.,

" $a_n = P(n)$  polyn of deg  $\leq k$ "

$\Leftrightarrow$

$(\Delta^m a)_0$  terminates for  $m > k$ .

"

$$\sum_{i=0}^m (-1)^i \binom{m}{i} a_{m-i}$$

Ex  $(0, \dots, 0, \frac{1}{k}, 0, \dots) \rightsquigarrow \binom{n}{k}$   
 Newton's family

① Classical Dold-Kan equivalence of cats

③

$$C^{\circ}: SAb \longleftrightarrow Ch_2(Ab) : N$$

$\parallel$   
 $Ab_{\Delta}$   
 smp abgp

where

$$C(A_0) = \frac{A_n}{deg.} \quad \text{"numbered chain"}$$

$$d = \sum_{i=0}^n (-1)^i d_i$$

$N$  a nerve construction: two steps

(1) universal cocylinder object

$$C^{\circ} : \Delta \longrightarrow Ch(Ab)$$

$\geq 0$

$$[n] \mapsto C(\mathbb{Z}\Delta^n)$$

Rmk  $Ab_{\Delta}$  present cat,  
 so gen by colms fr  $\mathbb{Z}\Delta^n$ ,  
 so dec by right adj is engh

(2)

$$N(B_{\bullet}) := \underline{Hom}(C^{\circ}, B_{\bullet}) \in Ab_{\Delta}$$

Ex

$$C^0 = \mathbb{Z}$$

$$C^1 = \mathbb{Z} \oplus \mathbb{Z} \longleftarrow \mathbb{Z}$$

$$C^2 = \mathbb{Z}^3 \longleftarrow \mathbb{Z}^2 \longleftarrow \mathbb{Z}$$

⋮

$$N(B_{\bullet}) = B_0 \longleftarrow \left\{ \begin{array}{l} b \in B_1 \\ b_0, b_1 \in B_0 \\ db = b_1 - b_0 \end{array} \right\} \longleftarrow \left\{ \begin{array}{l} \\ \\ B_2 \oplus B_1^2 \oplus B_0 \end{array} \right\}$$

$\parallel$   
 $B_1 \oplus B_0$

General formula:

$$N(\mathbb{B}_0)_n \cong \bigoplus_{[n] \twoheadrightarrow [k]} \mathbb{B}_k$$

surj.

↑ collect back degenerate cells

Pass to ranks: Get Newton's family:

$$\sum_{k=0}^n \binom{n}{k} b_k$$

ie,

$$\begin{array}{ccc}
 \text{Ab} \triangle & \xrightleftharpoons{\text{Dold-Kan}} & \text{Ch}_{20}(\text{Ab}) \\
 \text{rks} \downarrow & & \downarrow \text{rks} \\
 \{(a_n)\} & \xrightleftharpoons{\text{Newton}} & \{(b_n)\}
 \end{array}$$

(covered by Jergal)

Feature,  $|N(\mathbb{B}_n)| \triangleq K(\mathbb{B}_n)$ .

$$\binom{n}{k} \leftarrow (q \dots 0, 1, q \dots 0, 1)$$

② Secondary: dg cuts up to Morita equiv.

⑤

Note " $\sum (-1)^i d_i$ " doesn't make sense, since can't take differences of factors.  
dg cuts are com additive, not additive.

Fact  $\neq$  not quotient of  $A_n$ , but subobjects of  $A_n$  constrn of DK. (of  $C$ ).

The alternative constrn:

$$Ab_{\Delta} \ni A_0 \longmapsto C(A_n)_n = \bigcap_{\substack{i=1 \\ (\text{NOT } i=0)}}^k \text{Ker}(d_i)$$

$$d = d_0.$$

↑ makes sense for dg cuts.

Then  $\bigcap \text{Ker} \hookrightarrow A \twoheadrightarrow A/\text{deg}$  composition is an equivalence!

Ex  $C^1 = \mathbb{Z} \oplus \mathbb{Z} \leftarrow \mathbb{Z}$   
 $\{0, 1\} \quad \{0, 1, \infty\}$

$C^2 = \mathbb{Z} \oplus \mathbb{Z} \oplus \mathbb{Z} \leftarrow \mathbb{Z}^3 \twoheadrightarrow \mathbb{Z}$   
 $\{0, 1, 2\} \quad \begin{matrix} (0, 1, 0) \\ (0, 2, 0) \\ (2, 1, 1) \end{matrix} \quad (0, 1, 2 - 0, 1, 1 - 1, 0, 0, 1, 1)$

For  $A_0 \in (dg\text{-Cat})_{\Delta}$  i.e. from  $N(\mathcal{A}^{op}) \rightarrow dg\text{-Cat}$   
 as cat, recalled @ Mar 4 (6)

let

$$C(A) := \bigcap_{i=1}^n \text{Ker}(d_i)$$

make sure to talk of ker.

$$d = d_0$$

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{\text{Ker}} & \text{Part}(\mathcal{C}) \\ \downarrow & \sim & \downarrow \quad \downarrow \\ \mathcal{D} & & \mathcal{D} \end{array}$$

what maps

→ "chain complex of dg cat"

$$\text{Ch}_{\geq 0}(dg\text{-Cat})$$

$C(A) =$  "acyclic number of chains"

Secondly note:

$$\begin{aligned} \mathcal{C}^0: \mathcal{A} &\rightarrow \text{Ch}_{\geq 0}(dg\text{-Cat}) \\ [n] &\mapsto \underbrace{\mathbb{Z}\langle \text{Hom}(-, [n]) \rangle}_{\text{Cat}_{\Delta}} \end{aligned}$$

$$\underline{\text{Def}} \quad N(B_0) := \underline{\text{Hom}}_{\text{Cat}_{\Delta}}(\mathcal{C}^0, B_0).$$

$$C^1 = \mathbb{Z}\{0 \rightarrow 1\} \longleftarrow \mathbb{Z}\{01-00\}$$

↑  
"take care of"

↑  
"pairs"

= In cat gen by  
this  
=  $A_2$ -Mod.

↑  
 $A_1$ -Mod

$$1-0 \xleftarrow{d_0} (01-00)$$

↑  
Cone(0→1)

$$C^2 = \mathbb{Z}\{0 \rightarrow 1 \rightarrow 2\} \xleftarrow{d_0} \mathbb{Z} \left\{ \begin{array}{l} (01-00) \rightarrow (02-00) \\ (00-10) \rightarrow (02-00) \\ \downarrow \\ (12-11) \end{array} \right\} \xleftarrow{d_0} \mathbb{Z} \left\{ \begin{array}{l} (012-0011) \\ (002-001) \end{array} \right\}$$

=  $A_3$ -Mod

↑  
 $A_2$ -Mod

↑  
 $A_1$ -Mod

total of  
2-term  
complex

$K_0 \rightsquigarrow$  usual expts in k-obj.

$$\text{In gen: } (C^n)_{k-1} \cong \mathbb{Z} \left\{ \begin{array}{l} \text{part of } (k-1)\text{-slices} \\ \text{in } \Delta^n \end{array} \right\} / (\text{equiv obj}) \cong T(\Delta^k)$$

↑

$(C^n)_k$

↑

$\Delta^k$

$$T = \text{tot} (\partial_k \Delta^k \rightarrow \partial_{k-1} \Delta^k \rightarrow \dots \rightarrow \partial_0 \Delta^k)$$

## Features:

(8)

$N(BF_n) \cong K(B, n)$  "ordinary" EM spec.

•  $K(B, 1) = S_0(B)$  s-dot custom of  $B$ !

•  $K(B, 2) = S_0^{2,1}(B)$  Harish-Chandra-Macdonald  
"real s. custom"  
used for IR K-thy.

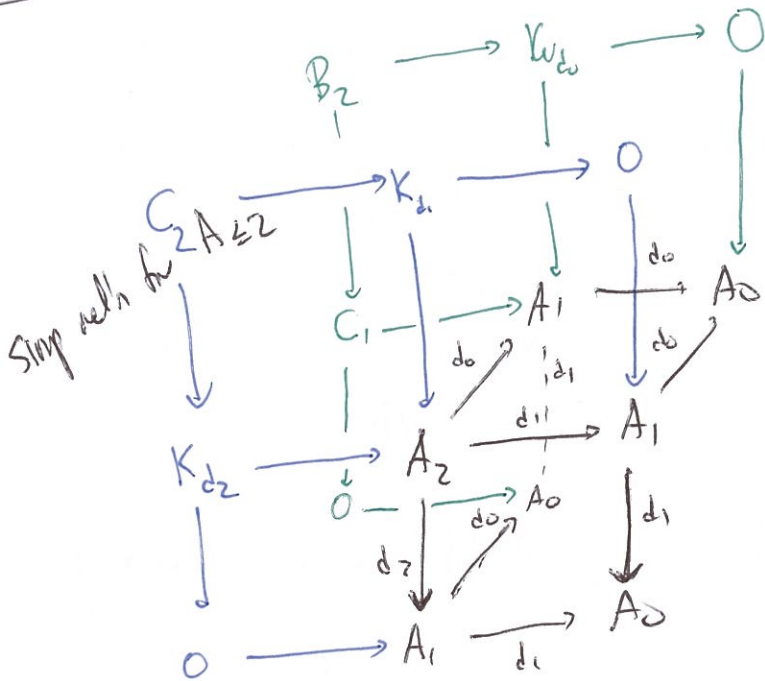
•  $K(B, n)$  could be called " $S_0^{n,1}(B)$ ", considered by G. Jasso  
for  $n$ -abelian cats

•  $N(B_1 \xrightarrow{f} B_2) \cong S_0(f)$  s-dot rel  $f$ .

•  $K(B, n)$  satisfy higher analogs of 2-Segal properties.  
odd  $n$ -Segal spaces



Dyckhoff talk scratchwork

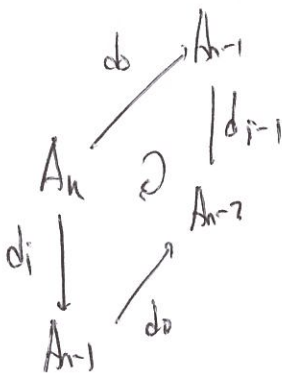


Pattern:

$$\left( \bigcup_{i=1}^n A_n \xrightarrow{d_i} A_{n-1} \right) \xrightarrow{d_0} \left( \bigcup_{j=0}^{n-1} A_{n-1} \xrightarrow{d_j} A_{n-2} \right)$$

$\mathbb{M}$   $\mathbb{M}$   
 $\text{Fun}(\bigvee_n \Delta', \text{Cat}_{\infty})$   $\text{Fun}(\bigvee_n \Delta', \text{Cat}_{\infty})$

$d_0 d_i = d_{i-1} d_0$



This induces map

$$\left( \bigcup_{i=1}^n A_n \xrightarrow{d_i} A_{n-1} \right) \xrightarrow{\begin{smallmatrix} d_0 \\ \circ \\ \circ \end{smallmatrix}} \left( \bigcup_{j=0}^{n-1} A_{n-1} \xrightarrow{d_j} A_{n-2} \right)$$

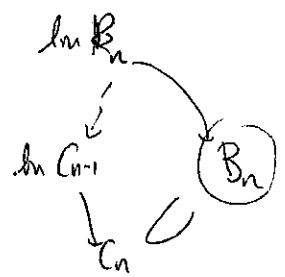
$\cap$   $\cap$

$$\text{Fun} \left( \underset{\underline{C}_n}{\bigvee_n \Delta'_1 \cup \Delta'_1}, \text{Cat}_\infty \right)$$

$$\text{Fun} \left( \underset{\underline{B}_n}{\bigvee_n \Delta'_1 \cup \Delta'_1}, \text{Cat}_\infty \right)$$

hence map

$$\begin{array}{ccc} C_n & \xrightarrow{\text{"do"}} & B_n \\ \parallel & & \parallel \\ \lim C_n & & \lim B_n \end{array}$$



We also have map  $B_n \rightarrow C_{n-1}$  because of inclusion

$$\left( \bigcup_{j=1}^{n-1} A_{n-1} \xrightarrow{d_j} A_{n-2} \right) \hookrightarrow \left( \bigcup_{j=0}^{n-1} A_{n-1} \xrightarrow{d_j} A_{n-2} \right) \quad \begin{array}{c} \text{Cat}_\infty \\ \uparrow \\ B_n \end{array}$$

more rigorously,  $\underline{C}_{n-1} = \alpha^* \underline{B}_n$  where  $\alpha: \left( \bigvee_{n-1} \Delta'_1 \cup \Delta'_1 \right) \hookrightarrow \left( \bigvee_n \Delta'_1 \cup \Delta'_1 \right)$

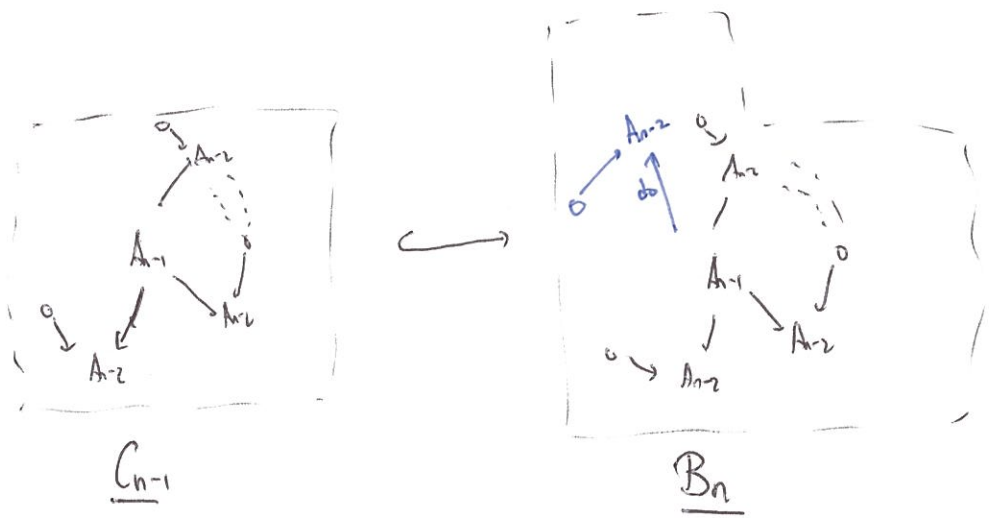
so have  $\lim B_n \rightarrow \lim \alpha^* C_{n-1}$ .

So the map

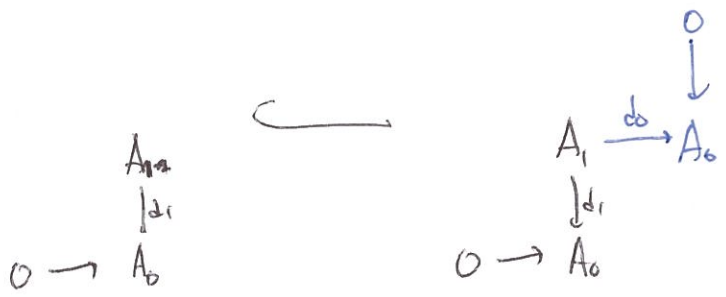
$$\begin{array}{ccccc}
 C_n & \xrightarrow{\text{"do"}} & B_n & \xrightarrow{\text{"d^*"}} & C_{n-1} \\
 & \searrow & & \nearrow & \\
 & & & \text{"do"} & 
 \end{array}$$

is what Toby calls do in the talk. Why does  $do \circ do \sim 0$ ?

Well, consider that



eg



so by stated pb propy,  $B_n \cong \text{p.b.} \left( C_{n-1} \xrightarrow{\text{do}} A_{n-1} \right)$ .

# Study composition

